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TITLE: INTERFACE LINK LAYER DEVICE TO BUILD A
DISTRIBUTED NETWORK

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Description

1 The present invention relates to an interface link layer device which builds the
core of an interface portal of an interface to which a network bus can be con-
nected for the purpose of building a distributed network by at least one other
network bus which is respectively connected to another interface portal of the
5 same interface.

Fig. 16 shows a coaxial interface between two IEEE1394 serial bus systems
which is existing according to the prior art and for example described in the
EP 0 848 568 A1 in a similar manner. An IEEE1394 serial bus 32 with three
10 IEEE1394 serial bus nodes 36 is connected to a first coaxial interface portal
31 which communicates via a coaxial interface, i. e. a coaxial cable 33, with a
second coaxial interface portal 34 to which a second IEEE1394 serial bus 35 is
connected which builds a network with two further IEEE1394 serial bus nodes
37. Each of the IEEE1394 serial bus nodes 36 has a different node identifier
15 as well as each of the further IEEE1394 serial bus nodes 37. Furtheron, each
of the first and second IEEE1394 serial busses 32 and 35 has a different bus
identifier. The node identifiers and the bus identifiers are assigned according
to the IEEE1394 standard. Therewith, the combination of node identifier and
bus identifier assigns a unique identifier to each of the serial bus nodes, even
20 if one of the serial bus nodes 36 and one of the further serial bus nodes 37 has
the same node identifier. Both IEEE1394 serial bus systems are independent
from each other, but they can communicate through the coaxial interface por-
tals and the coaxial interface.

25 Fig. 17 shows a multiportal system in which three interface portals 31, 34 and
38 are connected to one coaxial cable. Each of the interface portals 31, 34 and
38 is connected to an IEEE1394 serial bus having a different bus identifier.

As mentioned above, such systems are described in EP 0 848 568 A1 according
30 to which each interface portal is built by a RF/1394 converter that is
connected to a coaxial cable on one side and to an IEEE1394 bus on the other
side. Said RF/1394 converter converts all incoming data from the IEEE1394
bus into an RF-signal output to the coaxial cable and vice versa converts all
35 incoming RF-signals into data packets output to the connected IEEE1394
serial bus. To fulfill this task the RF/1394 converter basically comprises a
RF-modulator/demodulator, a link layer device and an IEEE1394 physical

1 layer device. Since in this case in the uplink channel basically all data packets
received on the IEEE1394 serial bus are just set into another format and
transmitted as RF-signal on a coaxial cable and vice versa in the downlink
channel all incoming RF-signals are converted into data packets output to the
5 IEEE1394 serial bus the link layer device has not to fulfill complicated tasks,
but it only secures the timing requirements for the IEEE 1394 serial bus.

However, according to this prior art many resources are wasted, since all data
packets that are available as RF-signals on the coaxial cable as well as all data
10 packets that are present on any one of the IEEE1394 serial busses are distrib-
uted within the whole network.

Therefore, it is the object of the present invention to provide an interface link
layer device that enables the interconnection of different networks to one dis-
15 tributed network while saving resources on said network.

Therefore, an interface link layer device according to the present invention that
is connected to a first data bus and via a transmission path to at least one
other interface link layer device that is respectively connected to a respective
20 second data bus is characterized by uplink means to accept a data packet hav-
ing a predetermined destination different to the interface link layer device it-
self from the first data bus and to transmit it via said transmission path to one
of said other interface link layer devices that is serving said predetermined
destination, and downlink means to output a data packet directly received via
25 said transmission path from one of said at least one other interface link layer
devices to a predetermined destination on the first data bus.

This interface link layer device according to the present invention is defined in
independent claim 1. Preferred embodiments thereof are respectively defined in
30 the dependent subclaims 2 to 33.

Therewith, according to the present invention every data bus that builds a part
of a distributed network is connected with an interface portal that comprises
as its core an interface link layer device according to the present invention
35 that directly communicates with all other interface link layer devices that are
working according to the present invention and that build respectively the core
of a respective interface portal respectively connected to a data bus building a
respective other part of said distributed network. With such a configuration

1 according to the present invention a data packet having a certain destination identifier is only delivered into such parts of the distributed network that comprise a receiver of this data packet. Such a distribution is preferably conducted on basis of the destination identifier, i. e. the destination identifier

5 comprises not only the node identifier of the respective destination node, but also the parts of the distributed network in which said respective node is located, e. g. a bus identifier of the network bus serving said part of the distributed network.

10 Therewith, this inventive method which is defined in independent claim 34 is applicable to a distributed network that is set up according to the present invention so that data channels are set up from every part of the distributed network to every other part thereof to allow a communication from one part of the distribution network to a predetermined other part thereof that is defined

15 in independent claim 35. A preferred embodiment of such a distributed network is defined in dependent claim 36.

Preferably, the present invention is used in an environment as described in the EP 0 848 568 A1, i. e. interconnects one or more separate IEEE1394 serial busses via coaxial interfaces, but it is also applicable to any other distributed network that is build through the interconnection of several individual networks respectively distributing data packets according to a predetermined standard to one distributed network. Of course, the individual networks need not all to work according to the same standard, i.e. the interface build with interface link layer devices according to the present invention can also interconnect networks working according to different standards.

The present invention and its embodiments will be better understood from a detailed description of an exemplary embodiment thereof taken in conjunction
30 with the accompanying drawings, wherein

Fig. 1 shows an interface link layer device according to a preferred embodiment of the present invention;

Fig. 2 shows the inserting of an IEEE1394 serial bus packet with less than
35 180 bytes into a 204 byte MPEG 2 frame;

Fig. 3 shows the separating of an IEEE1394 serial bus packet with less than 180 bytes from a 204 byte MPEG 2 frame;

Fig. 4 shows the fragmentation of a large serial bus packet before the trans-

1 mission over a coaxial cable;

5 **Fig. 5** shows the defragmentation of a large serial bus packet after its trans-
mission over the coaxial cable;

10 **Fig. 6** shows the acknowledge handling for asynchronous packets according
to a first example;

15 **Fig. 7** shows the acknowledge handling for asynchronous packets according
to a second example;

20 **Fig. 8** shows the acknowledge handling for asynchronous packets according
to a third example;

25 **Fig. 9** shows the acknowledge handling for asynchronous packets according
to a fourth example;

30 **Fig. 10** shows a simplified distributed network using a coaxial cable;

35 **Fig. 11** shows another simplified distributed network using a coaxial cable;

40 **Fig. 12** shows an additional function unit for an interface link layer device
according to the present invention;

45 **Fig. 13** shows a ring channel assignment of a distributed network with three
connected IEEE1394 serial busses;

50 **Fig. 14** shows a semi-hyper cube channel assignment of a distributed net-
work with three connected IEEE1394 serial busses;

55 **Fig. 15** shows a full-hyper cube channel assignment of a distributed network
with three connected IEEE1394 serial busses;

60 **Fig. 16** shows a distributed network with two IEEE1394 serial busses
connected via a coaxial interface according to the prior art; and

65 **Fig. 17** shows the interface part of a distributed network with three
IEEE1394 serial busses that are interconnected via a coaxial inter-
face according to the prior art.

30 Although the interface link layer device according to the present invention also builds the core of a respective interface portal that connects a network bus to an interface of a distributed network, for the sake of simplicity the respective physical layer devices interconnected inbetween the interface link layer device and the respective physical network bus and transmission path are omitted in the description.

35 Furtheron, the following description refers only to features that are not present in standard link layer devices. Of course, the interface link layer devices according to the present invention might also include the standard features, e.g. according to the IEEE1394 standard, that are not mentioned in the follow-

1 ing. For example, the interface link layer device according to the present in-
vention might be adapted to accept asynchronous packets with a destination
identifier that is equal to the node identifier of the NODE_IDS register within
its control and status registers as defined in the IEEE1394 specification, i.e.
5 asynchronous packets that are addressed to the interface link layer device it-
self, and might also be adapted to insert asynchronous packets on its
IEEE1394 bus, e. g. as response to such packets addressed to the interface
link layer device itself.

10 The exemplary embodiment described in the following shows the connection of
two or more IEEE1394 serial busses 2, 5 by a coaxial interface, i. e. by a coax-
ial cable 3. Such a system can for example be used in a home network environ-
ment in which individual IEEE1394 serial bus systems are installed in differ-
ent rooms of a house and are interconnected via existing coaxial cables so that
15 for example a video tape reproduced by a video recorder in one room of the
house can be watched on a television set in another room of said house, as it
is described in the EP 0 848 568 A1. In contrast to the system described in
this document according to the present invention the data packets carrying
said video film are only transmitted to a predetermined television set in a pre-
20 determined room, i. e. to a predetermined node within a predetermined part of
the distributed network, and not broadcasted through the whole network.

25 To facilitate this an interface link layer device according to a preferred embodi-
ment of the present invention which is shown in Fig. 1 comprises an uplink
section, a downlink section and an acknowledge/response section.

The uplink section accepts data packets having predetermined destinations
which are different to the interface link layer device 1 itself, i.e. the interface
30 portal comprising the interface link layer device 1, from a first data bus 2 and
to transmit those data packets via a transmission path 3, that might be a co-
axial cable, to predetermined other interface link layer devices 4 respectively
serving one of said predetermined destinations. To filter such data packets the
uplink section comprises a first register 12 which stores all allowed pre-
35 determined destinations, i. e. all destination identifiers of devices that can be
reached through the coaxial cable and that are connected to another part of
the distributed network, and a second register 15 which stores all data
channel numbers set-up inbetween two nodes of the distributed network, i.e.
inbetween two nodes located in different parts of the distributed network.

1 Preferably not the destination identifiers themselves, but only the bus identifier is stored and checked that defines a bus building one other part of the distributed network.

5 A control section 11 monitors all asynchronous data packets that are received from the IEEE1394 serial bus 2, compares their bus and/or node identifier with the bus and/or node identifiers stored in the first register 12 and opens or closes a switch 10 to lead an incoming asynchronous data packet to the further processing stages or to discard it depending on whether or not the destination identifier of the incoming packet matches with a destination identifier stored in the first register 12. Furtheron, the control section monitors all isochronous data packets that are received from the IEEE1394 serial bus 2, compares their channel number with the channel number stored in the second register 15 and opens or closes said switch 10 to lead an incoming isochronous data packet to the further processing stages or to discard it depending on whether or not the channel number of the incoming packet matches with a channel number stored in the second register 15.

20 The first register 12 is preferably a BUS_ENABLE register which is a bitmap that represents whether or not transaction requests and responses are forwarded by the interface. For example, if bit 2 is set to one in the BUS_ENABLE register, the interface link layer is enabled to accept asynchronous data packets with a bus identifier of 2. If this bit is not set, the interface link layer device denies the packet as described above. The second register 15 is preferably 25 a STREAM_CONTROL register which is similar organized as the first register 12 and has a corresponding function. Of course, as mentioned above, independently from this invention the interface link layer device always accepts data packets which destination identifier matches with its own node identifier, as every standard interface link layer device does.

30 The accepted IEEE1394 packets which should be transmitted via the coaxial cable to another part of the network are received by a packetizer 13 that repacks them into a data format suitable for the transmission channel, i. e. the transmission via the coaxial cable. Such a repacking is shown in detail in 35 the following in connection with Figs. 2 and 3. Fig. 2 shows the inserting of an IEEE1394 serial bus packet with a maximum of 180 bytes into a 204 byte frame according to the MPEG 2 standard. Since it is possible to transmit a data block of 180 bytes in the first packet of a section according to the MPEG

1 2 standard, the IEEE1394 serial bus packet including its header, payload and
cyclic redundancy code CRC is inserted as data block or part of the data block
into one 204 byte frame according to the MPEG 2 standard. If the whole
5 IEEE1394 serial bus packet comprises less than 180 bytes, i. e. less than the
data block of the MPEG 2 frame, some stuff bytes will be added to fill this dat-
ablock. Then the whole MPEG 2 frame including its header, its data block
which comprises the header, the payload and the cyclic redundancy code of
the IEEE1394 serial bus packet and some stuff bytes, as well as a 16 byte
Reed Solomon parity or another error detection/correction code is transmitted
10 via the coaxial cable.

If, on the other hand a large serial bus packet should be transmitted over the
coaxial cable 3 which has more than 180 bytes in total, a fragmentation has to
be performed before it is transmitted over the coaxial cable 3. Such a fragmen-
15 tation is shown in Fig. 3. Fig. 3 shows an example in which an IEEE1394 se-
rial bus packet of 732 bytes in total, i. e. including header, payload and cyclic
redundancy code is transmitted within four 204 byte frames according to the
MPEG 2 standard on the coaxial cable which each consists of a 8 byte header,
a 180 byte data block and a 16 byte Reed Solomon parity.

20 As can be seen in Fig. 1, the packetizer 13 is connected to said STREAM_

CONTROL register 15 which stores beside the channel number of each isochronous
channel the payload which defines the maximum number of quadlets that may
be transmitted in a single isochronous packet and which corresponds to the
25 data rate. This data is used by the packetizer 13 to properly set-up an isochro-
nous channel according to the MPEG 2 standard on the coaxial cable 3. Furtheron,
the packetizer 13 is connected to a third register, namely a BAND-
WIDTH_AVAILABLE register which provides the available bandwidth on the co-
axial cable 3 so that the packetizer 13 can check whether there is enough
30 bandwidth in case a new isochronous channel should be set-up or an asyn-
chronous packet should be transmitted.

A further element of the uplink section is a fourth register 17, namely an IN-
TERFACE_CONTROL register which includes the node identifier of the interface
35 link layer device connected to the other side of the distributed network, i.e. the
other side of the coaxial cable 3, and an Enable Bit to allow the direct routing
of asynchronous packets through the coaxial cable 3.

- 1 The data packets to be output on the downlink channel(s) of the IEEE1394 serial bus 2 through the interface link layer device 1 are received through the connected transmission channel 3 from another interface link layer device 4 by the downlink section of the interface link layer device 1. Within said downlink
- 5 section said data packets get repacked by a packet separator 20 into the IEEE1394 format which has a reverse operation in comparison with the packetizer 13 of the uplink section.

As it is shown in Fig. 4 for the case that an IEEE1394 serial bus packet with
10 less than 188 bytes gets separated from an incoming MPEG 2 frame which has 204 bytes the packet separator 20 takes the header, the payload and the cyclic redundancy code which belong to the IEEE1394 serial bus packet from the data block of the MPEG 2 frame and discards the stuff bytes which have been added by the packetizer 13 of the other interface link layer device 4. The case
15 that the IEEE1394 serial bus packet has a length of more than 188 bytes and needs a defragmentation after transmitting over the coaxial cable is shown in Fig. 5. In the shown case the IEEE1394 serial bus packet has a total length of 752 bytes and is therefore transmitted in the data blocks of 4 MPEG 2 frames. The data blocks of these 4 frames carry the header, the payload and the cyclic
20 redundancy code of the IEEE1394 serial bus packet which gets assembled in the packet separator 20 and distributed via the downlink channel(s) of the IEEE1394 serial bus 2.

Since the packetizer 13 adds stuff bytes to fill each data block of the 204 byte
25 MPEG 2 frame and in case that there are no IEEE1394 packets to transmit via the coaxial cable a data block of the 204 byte MPEG frame consists only of stuff bytes the packet separator 20 receives a constant data stream from which the IEEE1394 packets get separated. As described above, divided IEEE1394 packets will be put together again and then, if not addressed to the outbound
30 portal, i. e. the interface link layer device 1 itself, will be transmitted by the downlink section within the interface link layer device 1 to the connected IEEE1394 serial bus 2 to the destination node.

If a predetermined bit in the header of a frame indicates a transport error the
35 interface link layer device discards the complete packet. Also, if an error occurs during transmission via the coaxial cable the packet will be discarded. In these cases a device which waits for said packet, since it has requested a certain transaction, generates a split timeout error which is an indicator that the

1 certain transaction should be requested again. The minimum timeout value for a split timeout is 100 ms.

As described above, in case of asynchronous packets the header is necessary
5 to identify the source, the destination and the transaction type of the packet.
For isochronous packets, on the other hand, there is no guarantee to be trans-
mitted within the same channel number from the outbound portal. Therefore,
the downlink section comprises a channel number assignment unit 21 that as-
signs a new channel number to isochronous packets which origin in a channel
10 of the data bus that hosts the originator of said packets which channel num-
ber is already occupied for another data stream within the IEEE1394 serial
bus 2 that is connected to the interface link layer device 1 and serves the re-
ceiver of said packets.

15 As described hereinafter in connection with Figs. 6 to 9 the interface link layer device according to the present invention does not just forward incoming acknowledge packets to their destination, but sets up an own acknowledgement and response strategy. Therefore, acknowledge packets that are received via the IEEE1394 serial bus 2 are not just forwarded through the coaxial cable 3
20 to the interface link layer device connected to the network bus to which the destination device is connected, but the incoming acknowledge packet gets analysed by said interface link layer device which then decides which kind of action, i. e. acknowledgement and/or response is necessary to which device.

25 To properly set-up an acknowledge and response strategy the acknowledge/response section of the interface link layer device 1 according to the present invention comprises an acknowledge code generator 18 that serves the inbound behaviour and a response packet generator 19 that serves the outbound behaviour.

30 The acknowledge code generator 18 receives all accepted data packets through the uplink channel(s) of the IEEE1394 serial bus 2 and generates appropriate acknowledge codes to be output on the downlink channel(s) of the IEEE1394 serial bus 2. Therefore, according to the present invention, an acknowledge
35 code is not only generated when a packet which is to be transmitted to another part of the distributed network, i. e. another IEEE1394 serial bus, through the coaxial cable 3 reaches its destination, but also when such a packet reaches the inbound portal, i. e. the interface link layer device 1, of the trans-

1 mission channel, i. e. the coaxial cable 3.

The response packet generator 19 that monitors the uplink and downlink channel(s) of the connected IEEE1394 serial bus 2 generates an appropriate 5 response packet to be output with an appropriate speed to a particular destination device via the coaxial cable 3 in case a request was transmitted to the connected IEEE1394 serial bus 2 and a corresponding acknowledge code is received from the connected IEEE1394 serial bus 2.

10 As mentioned above, the acknowledge code generator 18 sends an acknowledge to the originator of a packet received via the connected IEEE1394 serial bus 2 for each accepted asynchronous packet that is either addressed to the node controller of the interface link layer device 1 or to a node on the other side of the coaxial cable 3, e.g. a node connected to another IEEE1394 serial bus 5 15 within the distributed network. The following table 1 gives an overview about the outbound behaviour of the interface link layer device concerning the acknowledge. A "-" sign within the table indicates that the respective content does not matter, but differs from 3F or 3FF. 3F and 3FF are hexadecimal values which are reserved within the IEEE 1394 standard for special purposes.

15 The Enable Bit is located in the INTERFACE_CONTROL register 17, as stated above.

Enable Bit	Destination Bus ID	PHY ID	Source Bus ID	Action
25	3FF or Node_IDS.bus_ID	-	-	Local packet. Acknowledge per IEEE1394-1995 standard.
	Neither 3FF nor Node_IDS.bus_ID	-	-	Routing disabled. Ignore packet, no Acknowledge.
	-	-	3FF	Unable to route. Ignore packet, no acknowledge.
30	1	3FF	3F	Global broadcast. Forward packet, no acknowledge.
	1	3FF	Not 3FF	Remote packet origination. Forward packet according to BUS_ENABLE, transmit ack_pending or ack_complete.
35	1	Neither 3FF nor Node_IDS.bus_ID	Node_IDS.bus_ID	Remote packet in transit. Forward packet and transmit ack_pending or ack_complete.
			Neither 3FF nor Node_IDS.bus_ID	Unable to route. Ignore packet, no acknowledge
All other combinations				Unable to route. Ignore packet, no acknowledge

1 It can be seen that the generated acknowledge code depends on the destination identifier, the transaction code and the cyclic redundancy code of an incoming packet. If a received packet is a request and to be routed through the interface, the acknowledge code generator 18 always generates an acknowledge

5 code that indicates a pending action, e. g. ack_pending, for the indicator of the request. In case of an error the acknowledge code generator 18 generates a data error acknowledgement, e. g. ack_data_error, and an acknowledgement that indicates a completed action will be generate in case the received packet is a response packet without any errors, e. g. ack_complete. This behaviour of

10 split transactions is necessary, since in most cases there would be no unified transaction possible over interfaces for time reasons, because the originator of a packet expects an acknowledge code for this packet within a certain time limit, e.g. 50 ms.

15 The outbound behaviour, on the other hand, is controlled by the response packet generator 19. Every time an asynchronous packet, i.e. a request packet, was transmitted through the coaxial cable to a predetermined destination on the connected IEEE1394 serial bus the incoming corresponding acknowledge packet is received by the response packet generator 19 from the connected

20 IEEE1394 serial bus 2 and it is decided whether or not an appropriate response is output via the coaxial cable 3 to the originator of a request that has triggered the acknowledge, as shown in the following table 2.

	Name	Action
25	(incoming acknowledgement)	
	ack_complete	Request: Transmit the appropriate response packet. Response: No action.
	ack_pending	
	ack_busy_X	The bridge abandons retry attempts:
	ack_busy_A	Request: Transmit the appropriate response packet.
	ack_busy_B	Response: No action.
30	ack_data_error	Request: Transmit the appropriate response packet.
	ack_type_error	Response: No action.

Therefore, the response packet generator 19 monitors all accepted acknowledge packets that are incoming via the connected IEEE1394 serial bus 2 to provide an appropriate response packet as an asynchronous response packet for the originator of the request packet. Furtheron, the response packet generator 19 has to monitor all request packets output to the connected IEEE1394

1 serial bus 2. To monitor all incoming and outgoing packets the response packet generator 19 stores different components of the asynchronous packet header, e.g. the source and destination of a packet, the source as new destination and the destination as new source.

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It can be seen that the appropriate response packet will only be generated by the response packet generator 19 if the asynchronous packet transmitted via the coaxial cable 3 to a predetermined destination on the connected IEEE1394 serial bus 2 was a write request and the response packet generator 19 received

10 an acknowledge that indicates a completed action from the predetermined destination, i.e. the receiver of said request. For read requests and lock requests the receiver of the asynchronous packet has to create the appropriate response packet which will be forwarded by the other interface link layer device 4 through the coaxial cable 3 to the interface link layer device that is connected
15 to the IEEE1394 serial bus serving the originator of the corresponding request packet.

Figs. 6 to 9 respectively show an example in which a first IEEE1394 node 6 is connected via a first IEEE1394 serial bus 2 to a first interface link layer device
20 1 according to the present invention which is connected through a coaxial cable 3 to a second interface link layer device 4 according to the present invention that is connected to a second IEEE1394 serial bus 5 which serves a second IEEE1394 node 7.

25 In the first example shown in Fig. 6, the first node 6 distributes in a first step S61 a request to the second node 7 on the first serial bus 2 it is connected to. Such a request is send as an asynchronous packet. In a second step S62 the first interface link layer device 1 receives that request and accepts it, since the identifier of the second node 7 is stored in its first register 12, i. e. the
30 BUS_ENABLE register. Additionally, the acknowledge code generator 18 of the first interface link layer device 1 transmits an acknowledgement indicating a pending action, i.e. ack_pending, to the first node 6 and transfers the request via the coaxial cable 3 to the second interface link layer device 4 according to the MPEG 2 standard. In a third step S63 the second interface link layer device 4 receives the request, repacks it with its packet separator 20 and transmits it via the connected second serial bus 5 to the second node 7. The second node 7 indicates in a fourth step S64 a completed action with an acknowledge-

1 ment, i.e. ack_complete, directed to the requesting first node 6. This acknowledgement code is distributed on the second serial bus 5 which is connected to the second node 7. In the next fifth step S65 the second interface link layer device 4 receives and accepts the acknowledge packet from the second node 7,
5 since its destination identifier is stored in its first register 12, i.e. its BUS_ENABLE register. The response packet generator 19 of the second interface link layer device 4 generates an appropriate response packet and transmits it via the coaxial cable 3 to the first interface link layer device 1 which distributes it after an appropriate repacking with its packet separator 20 to the first node 6
10 via the first serial bus 2 in a sixth step S66. Finally, in the seventh step S67, the first node 6 sends an completed action acknowledgement code to the second node 7 via the connected first serial bus 2 which is received and accepted by the first interface link layer device 1 which accepts and discards it and therewith finalizes the action.

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In the second example shown in Fig. 7, the first node 6 distributes in a first step S71 a request to the second node 7 on the first serial bus 2 it is connected to as an asynchronous packet. In a second step S72 the first interface link layer device 1 receives that request and accepts it, since the identifier of

20 the second node 7 is stored in its first register 12, i. e. the BUS_ENABLE register. Additionally, the acknowledge_code_generator_18 of the first interface link layer device 1 transmits an acknowledgement indicating a pending action, i.e. ack_pending, to the first node 6 and tranfers the request via the coaxial cable 3 to the second interface link layer device 4 according to the MPEG 2 standard. In a third step S63 the second interface link layer device 4 receives the request, repacks it with its packet separator 20 and transmits it via the connected second serial bus 5 to the second node 7. The second node 7 indicates in a fourth step S74 a pending action with an acknowledgement, i.e. ack_pending, directed to the requesting first node 6. This acknowledgement code is distributed on the second serial bus 5 which is connected to the second node 7. Therafter, in a sixth step S76, the second node 7 generates a response to the request received in the third step S73 that is directed to the requesting first node 6. This response is distributed on the second serial bus 5 which is connected to the second node 7 and received and accepted by the second interface link layer device 4, since its destination identifier is stored in its first register 12, i.e. its BUS_ENABLE register. The acknowledge code generator 18 of the second interface link layer device 4 transmits an acknowledge code indicating

1 a completed action, i.e ack_complete, to the second node, since the second interface link layer device has received the response without any errors. Since the response, which indicates a completed action, is no acknowledge code the second interface link layer device 4 transmits said packet after repacking by

5 its packetizer 13 via the coaxial cable 3 to the first interface link layer device 1 which distributes it after an appropriate repacking with its packet separator 20 to the first node 6 via the first serial bus 2 in a seventh step S77. Finally, in the eighth step S78, the first node 6 sends an completed action acknowledgement code to the second node 7 via the connected first serial bus 2 which

10 is received and accepted by the first interface link layer device 1 which accepts and discards it and therewith finalizes the action.

The timeout value of the first node 6 to wait from the received ack_pending acknowledgement to the response should be set >100 ms.

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In the third example shown in Fig. 8, the first node 6 distributes in a first step S81 a request to the second node 7 on the first serial bus 2 it is connected to as an asynchronous packet. In a second step S82 the first interface link layer device 1 receives that request and accepts it, since the identifier of the second

20 node 7 is stored in its first register 12, i. e. the BUS_ENABLE register. Additionally, the acknowledge code generator 18 of the first interface link layer device 1 transmits an acknowledgement indicating a pending action, i.e. ack_pending, to the first node 6 and transfers the request via the coaxial cable 3 to the second interface link layer device 4 according to the MPEG 2 standard.

25 In a third step S83 the second interface link layer device 4 receives the request, repacks it with its packet separator 20 and transmits it via the connected second serial bus 5 to the second node 7. The second node 7 indicates in a fourth step S84 that it is busy with an acknowledgement, i.e. ack_busy, directed to the requesting first node 6. This acknowledgement code is distributed

30 on the second serial bus 5 which is connected to the second node 7. In the next fifth step S85 the second interface link layer device 4 receives and accepts the acknowledge packet from the second node 7, since its destination identifier is stored in its first register 12, i.e. its BUS_ENABLE register. The response packet generator 19 of the second interface link layer device 4 generates an appropriate response packet and transmits it via the coaxial cable 3 to the first interface link layer device 1 which distributes it after an appropriate repacking with its packet separator 20 to the first node 6 via the first serial

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1 bus 2 in a sixth step S86. Finally, in the seventh step S87, the first node 6
5 sends an completed action acknowledgement code to the second node 7 via the
connected first serial bus 2 which is received and accepted by the first inter-
face link layer device 1 which accepts and discards it and therewith finalizes
the action.

In the fourth example shown in Fig. 9, the first node 6 distributes in a first
10 step S91 a request to the second node 7 on the first serial bus 2 it is con-
nected to as an asynchronous packet. In a second step S92 the first interface
link layer device 1 receives that request and accepts it, since the identifier of
the second node 7 is stored in its first register 12, i. e. the BUS_ENABLE regis-
ter. Additionally, the acknowledge code generator 18 of the first interface link
layer device 1 transmits an acknowledgement indicating a pending action, i.e.
ack_pending, to the first node 6 and tranfers the request via the coaxial cable
15 3 to the second interface link layer device 4 according to the MPEG 2 stan-
dard. In a third step S83 the second interface link layer device 4 receives the
request, repacks it with its packet separator 20 and transmits it via the con-
nected second serial bus 5 to the second node 7. The second node 7 indicates
in a fourth step S84 a data error with an acknowledge code, i.e. ack_data_er-
ror, directed to the requesting first node 6. This acknowledgement code is dis-
tributed on the second serial bus 5 which is connected to the second node. In
the next fifth step S85 the second interface link layer device 4 receives and ac-
cepts the acknowledge packet from the second node 7, since its destination
20 identifier is stored in its first register 12, i.e. its BUS_ENABLE register. The re-
sponse packet generator 19 of the second interface link layer device 4 gener-
ates an appropriate response packet and transmits it via the coaxial cable 3 to
the first interface link layer device 1 which distributes it after an appropriate
repacking with its packet separator 20 to the first node 6 via the first serial
bus 2 in a sixth step S86. Finally, in the seventh step S87, the first node 6
30 sends an completed action acknowledgement code to the second node 7 via the
connected first serial bus 2 which is received and accepted by the first inter-
face link layer device 1 which accepts and discards it and therewith finalizes
the action.

35 It can be seen that according to the present invention the interface link layer
device which is connected to a respective IEEE1394 serial bus and a respective
coaxial cable through respective physical layer devices provides accepted pack-

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1 sets that should be transmitted via the interface as complete packets to the co-
axial cable. Such a complete packet consists of the header, the data block and
the cyclic redundancy codes. The transmission of the complete packet through
the coaxial cable 3 is necessary, because the corresponding interface link layer
5 device connected to another network bus of the distributed network needs the
information of the packets for the transmission on said other serial bus. Fur-
theron, the interface link layer device according to the present invention intro-
duced the concept that not all packets have to be transmitted via the interface,
since an appropriate acknowledge code and response code generation of the in-
10 terface link layer device itself is introduced.

Since according to the present invention not all, but only selected data packets
are distributed via the coaxial interface it is in a special arrangement possible
to "hide" the coaxial interface to the IEEE1394 bus to which it is connected.

15 An example for such an arrangement is shown in Fig. 10.

Generally, such arrangements comprise one first IEEE1394 serial bus 2 which
can have any number of controller nodes 6 connected. This is called the multi-
device side. In the example shown in Fig. 10 there are two IEEE1394 serial bus
20 nodes 6 and one interface link layer device 1 (which is also a node within the
IEEE1394 system) connected to the first IEEE1394 serial bus 2.

Via a coaxial cable 3 and a second interface link layer device 4 there is con-
nected a second IEEE1394 serial bus 5. According to this embodiment of the
25 invention to "hide" the coaxial interface 3 to the first IEEE1394 bus 2 this
second IEEE1394 serial bus 5 is allowed to have only one single IEEE1394 se-
rial bus node 7 connected. This is called the single-device side.

To "hide" the coaxial interface 3 it is necessary that a controller node 6 on the
30 multi-device side is not aware of the interface. This can be achieved since the
interface link layer device 1 looks like the single IEEE1394 serial bus node 7
connected to the first interface link layer device 1 via the coaxial interface 3,
the second interface link layer device 5 and the second IEEE1394 serial bus 5.

35 Fig. 10 shows also the node identifiers of all serial bus nodes connected to
both IEEE1394 serial busses 2 and 5 which are automatically assigned after a
bus reset. The multi-device side comprises two controller nodes 6, one has a

1 node ID 0 and the other has a node ID 2. The first interface link layer device 1
connected to the first IEEE1394 serial bus 2 has a node ID 1. The single-de-
vice side is build by a controller node 7 having a node ID 0 and the second in-
terface link layer device 4 having a node ID 1.

5

In order to make the communication work the destination node ID within asyn-
chronous packets of a controller node 6 connected to the first IEEE1394 serial
bus 2 to the single-device side has to be changed from the node ID of the first
interface link layer device 1 on the multi-device side, 1 in Fig. 10, to the node
10 ID of the single controller node 7 on the single device side, 0 in Fig. 10. The
source node ID remains unchanged for this communication direction.

Further, in asynchronous packets from the single-device side to the multi-de-
vice side the source node ID has to be changed from the node ID of the single
15 controller node 7, 0 in Fig. 10, to the node ID of the interface link layer device
1 on the multi-device side, 1 in Fig. 10. The destination ID remains unchanged
for this communication direction.

A possible implementation can be realized by providing two additional registers

20 within the second interface link layer device 4 of the single-device side. The
first-additional-register-stores-the-node-ID-of-the-first-interface-link-layer de-
vice 1 on the multi-device side and the second additional register stores the
node ID of the single controller node 7 on the single-device side. The second
interface link layer device 4 has then to replace the destination node identifier
25 before sending a packet to the second IEEE1394 serial bus 5 using the value
of the second additional register and to replace the source node identifier us-
ing the value of the first register before sending a packet to the coaxial cable 3.
In this implementation the first interface link layer device 1 connected to the
multi-device side requires only the modification to forward packets directed to
30 itself through the coaxial cable to the second interface link layer device, i.e. to
forward packets with destination node identifiers that are identical to its own
node identifier. Furthermore, the second interface link layer device 4 on the
single-device side has to operate in a special mode, because it has to accept
and forward packets with destination node identifiers that are not identical to
35 its own node identifier.

Of course, this embodiment to "hide" the coaxial interface is not limited to a

1 distributed network with only one pair of coaxial interfaces. Fig. 11 shows an example where two separate single remote nodes 7A, 7B are connected to a coaxial cable network 3. To achieve such a "hiding" of two separate single remote nodes two first interface link layer devices 1A, 1B are required on the multi-
5 device side, i. e. connected to the first IEEE1394 serial bus 2, which respectively communicate with a corresponding second interface link layer device 4A, 4B through the coaxial cable 3 which is in turn connected to a respective single remote node 7A, 7B through a respective second IEEE1394 serial bus 5A, 5B.

10

As described in EP 0 848 568 A1 there are different solutions for the transmission of digital data via the coaxial cable, e. g. as a QPSK or a QAM-signal, of course, all these and other methods can be applied to an interface link layer device according to the present invention.

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Figs. 13 to 15 show that more than two interface link layer devices can be connected to the coaxial cable 3 similar to the distributed prior art network shown in Fig. 17. It is shown that three interface link layer devices are connected to one coaxial cable 3, namely a first interface link layer device 1, a second inter-
20 face link layer device 4 and a third interface link layer device 8. Every interface link layer device is also connected to a respective IEEE1394 serial bus and each of these IEEE1394 serial busses has been assigned a different bus identifier.

25 According to the present invention a bi-directional communication between all busses is enabled and data that is not needed on a certain IEEE1394 serial bus will not be seen there. As stated above within the coaxial cable or a coaxial network consisting of several coaxial cables channels are assigned in frequency division multiplex with every channel occupying a certain bandwidth so
30 that a large number of channels is simultaneously available.

In order to achieve availability of the data packets within the coaxial cable 3 to all interface link layer devices while not setting-up two channels between every possible pair of interface link layer devices, i. e. while not setting-up a full-hyper-cube channel assignment as shown and described in connection with Fig. 15, but setting-up e. g. a ring channel assignment, while not generating superfluous bus traffic on the attached IEEE1394 serial busses, an additional func-

1 tion unit as shown in Fig. 12 and described in the following can be imple-
2 mented inside an interface link layer device according to the present invention.

3 The downstream path from the coaxial cable 3 to the interface link layer device
4 and the upstream path from the interface link layer device to the coaxial cable
5 3 can be connected if necessary by a first additional switch 16a which is con-
6 trolled by the controller 11 which is additionally connected to the downstream
7 path. Asynchronous IEEE1394 packets taking this route will not be seen on
8 the IEEE1394 serial bus this interface link layer device is connected to, since
9 a second additional switch 16b which is also controlled by the controller 11
10 and which connects the downstream path with the packet separator 20 is open
11 in this case.

12 The controller 11 monitors the incoming downstream packets to control the
13 first and second additional switches 16a and 16b according to the destination
14 address of the packet. The switch position shown in Fig. 12 in which the first
15 additional switch 16a is open and the second additional switch 16b is closed
16 indicates the case when a packet is destined to the IEEE1394 serial bus this
17 interface link layer device is a member of. As described above, the first addi-
18 tional switch 16a will close and the second additional switch 16b will open for
19 packets destined to other bus numbers.

20 In case of isochronous packets, it is possible to send these both to the
21 IEEE1394 serial bus the interface link layer device is connected to and
22 through the upstream path to the next or other interface link layer devices on
23 the coaxial cable 3. In this case both additional switches 16a and 16b are
24 closed.

25 Of course, such a functionality can alternatively also be included in the packet
26 separator 20 which needs then a connection to the upstream path, e.g. directly
27 or via the packetizer 13, and in order to know how to distribute the received
28 packets also to the controller 11.

29 Fig. 13 shows a ring channel assignment for the distributed IEEE1394 serial
30 bus network according to which the data links between all the interface link
31 layer devices connected to the coaxial cable 3 are arranged in a ring. For the
32 ring channel assignment every interface requires one transmission and one re-

1 ception channel. In Fig. 13 the first interface link layer device 1 transmits on
channel 1 and receives on channel 3, the second interface link layer device 4
transmits on channel 2 and receives on channel 1 and the third interface link
layer device 8 transmits on channel 3 and receives on channel 2. For n inter-
5 faces n channel on the coaxial cable are required. With this channel assign-
ment there is no need to re-configure the data links on the coaxial cable 3 on
the fly.

10 Fig. 14 shows a semi-hyper-cube channel assignment for the distributed
IEEE1394 serial bus network which assures a data link from every interface
link layer device to every other interface link layer device. To save channels
every interface link layer device transmits on just one channel via the coaxial
cable 3 and reception is done on all the other channels used for transmission
by the other interface link layer devices. In Fig. 14 the first interface link layer
15 device 1 transmits on channel 1 and receives on channels 2 and 3, the second
interface link layer device 4 transmits on channel 2 and receives on channels 1
and 3 and the third interface link layer device 8 transmits on channel 3 and
receives on channels 1 and 2.

20 For such a semi-hyper-cube channel assignment every interface link layer de-
vice requires one transmission and $(n-1)$ reception channels in case of n inter-
face link layer devices. Therefore, n channels on the coaxial network are re-
quired. If said data links are desired, one interface link layer device must have
25 $(n-1)$ front end modules. Therefore, it is desirable to have a dynamic channel
assignment scheme.

30 Fig. 15 shows the full-hyper-cube channel assignment for the distributed
IEEE1394 serial bus network which assures a data link from every interface
link layer device to every other interface link layer device. In the full-hyper-
cube channel assignment one transmission and one reception channel are set
up inbetween each two interface link layer devices. Therefore, the additional
functionality of the interface link layer device 1 as shown in Fig. 12 and de-
scribed in connection therewith is not necessary. In Fig. 15 the first interface
link layer device 1 transmits on channel 1 to the second interface link layer
35 device 4 and on channel 2 to the third interface link layer device 8, and re-
ceives on channel 3 from the second interface link layer device 4 and on chan-
nel 5 from the third interface link layer device 8. The second interface link

- 1 layer device 4 transmits on channel 3 to the first interface link layer device 1 and on channel 4 to the third interface link layer device 8, and receives on channel 1 from the first interface link layer device 1 and on channel 6 from the third interface link layer device 8. The third interface link layer device 8 trans-
- 5 mits on channel 5 to the first interface link layer device 1 and on channel 6 to the second interface link layer device 4, and receives on channel 2 from the first interface link layer device and on channel 4 from the second interface link layer device 4.
- 10 In a distributed network with n interface link layer devices every interface link layer device requires $(n-1)$ transmission and $(n-1)$ reception channels. Therefore, for n interface link layer devices $n \cdot (n-1)$ channels on the coaxial cable 3 are required.

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